

# Lightweight Probability Theory for Verification

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## Motivation

The Miller-Rabin primality test takes a number  $n$  and returns either PRIME or COMPOSITE. If  $n$  actually is prime then it is guaranteed to return PRIME, and if  $n$  is composite then it will return COMPOSITE with **probability** at least one half. Successive calls are **independent**, so if  $n$  is composite then  $s$  consecutive results of PRIME will occur with **probability** at most  $2^{-s}$ .

**How can we specify and verify such an algorithm?**

To answer this question, we have created the following two theories in HOL:

- A language for expressing probabilistic programs.
- A formalization of (basic) probability theory.

## A Language for Probabilistic Algorithms

The programming language we use is the language of higher-order logic functions.

We define a type  $\mathbb{B}^\infty$  of infinite boolean sequences, and model a probabilistic function

$$f : \alpha \rightarrow \beta$$

with a corresponding deterministic function

$$F : \alpha \rightarrow \mathbb{B}^\infty \rightarrow \beta \times \mathbb{B}^\infty$$

This method of ‘passing around the random-number generator’ is also used in pure functional languages such as Haskell, and allows an elegant formulation of probabilistic programs in terms of state transforming monads.

## Formalizing Probability Theory

We build upon Harrison's construction of the real numbers; adding ingredients from mathematical measure theory to allow the essential concepts of **probability** and **independence** to be defined. This results in a lightweight probability theory.

This leads to an important result:

**Thm:** For all probabilistic programs constructed using our monadic primitives (including Haskell probabilistic programs), the returned value is independent of the returned sequence.

Note: the converse is not true:  $\lambda s.(s_0 = s_1, \text{stl } s)$

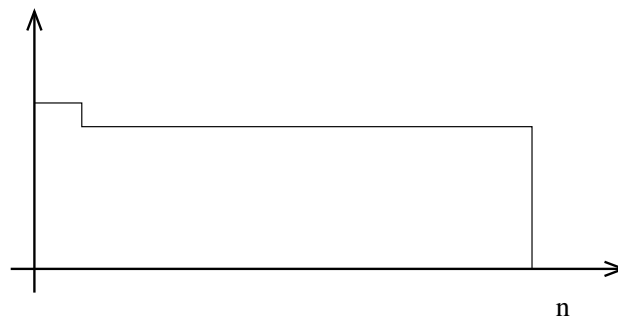
This indicates how tricky independence can be.

## A Uniform Random Number Generator

We made use of this development to write a probabilistic function that returned random numbers in the range  $0, 1, \dots, n - 1$ .

We originally wanted the returned numbers to be uniformly distributed on the range, but this turns out to be impossible unless  $n$  is a power of two!

We settled for almost-uniform:



We pass in an extra parameter  $t$ , and the probability of returning each number in the range is within  $2^{-t}$  of  $1/n$ .